CSCE 463/612 Networks and Distributed Processing Spring 2024

Introduction

Dmitri Loguinov Texas A&M University

January 24, 2024

Updates

- Recv loop reminder
 - timeout.tv_usec must be initialized to zero
 - NULL-terminate buf before searching with strchr or strstr

```
while (true) {
 if ((ret = select (0, &fd, ..., &timeout)) > 0)
   // new data available; now read the next segment
   int bytes = recv (sock, buf + curPos, allocatedSize - curPos, ...);
   if (errors)
        // print WSAGetLastError() & return false;
   if (connection closed)
                                                     commonly
       buf[curPos] = NULL;
        return true; // normal completion
                                                      forgotten
   if (allocatedSize - curPos < THRESHOLD)
        // realloc() buf to double its size
 else if (timeout)
        // report timeout & return false;
 else
        // print WSAGetLastError() & return false;
```

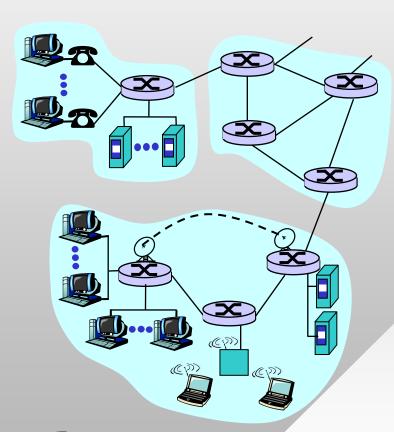
Chapter 1: Roadmap

- 1.1 What is the Internet?
- 1.2 Network edge
- 1.3 Network core
- 1.4 Network access and physical media
- 1.5 Internet structure and ISPs
- 1.6 Delay & loss in packet-switched networks
- 1.7 Protocol layers, service models
- 1.8 History

The Internet: "Nuts and Bolts" View

- 1) Hosts (end systems)
 - Computing devices (servers, desktops, phones, laptops)
 - Run network apps
- 2) Routers
 - Forward packets (chunks of data) to destinations
- 3) Communication links
 - Connect hosts & routers
 - Fiber, copper, radio, satellite
 - Transmissionrate = bandwidth

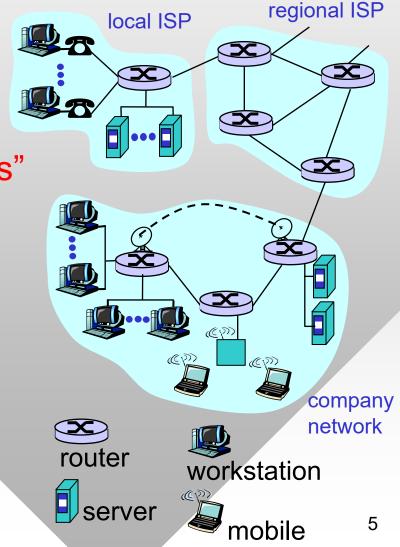






The Internet: "Nuts and Bolts" View

- 4) Protocols
 - Control sending/receiving of messages (e.g., TCP, IP, HTTP, FTP, SMTP)
- Internet: "network of networks"
 - Loosely hierarchical
- Who rules the Internet?
 - No single authority, mostly decentralized
- Internet standards
 - IETF: Internet Engineering
 Task Force
 - RFC: Request for comments



What's a Protocol?

Human protocols:

- "What's the time?"
- "I have a question"
- Introductions
- ... specific msgs sent
- ... specific actions taken when msgs received or other events take place

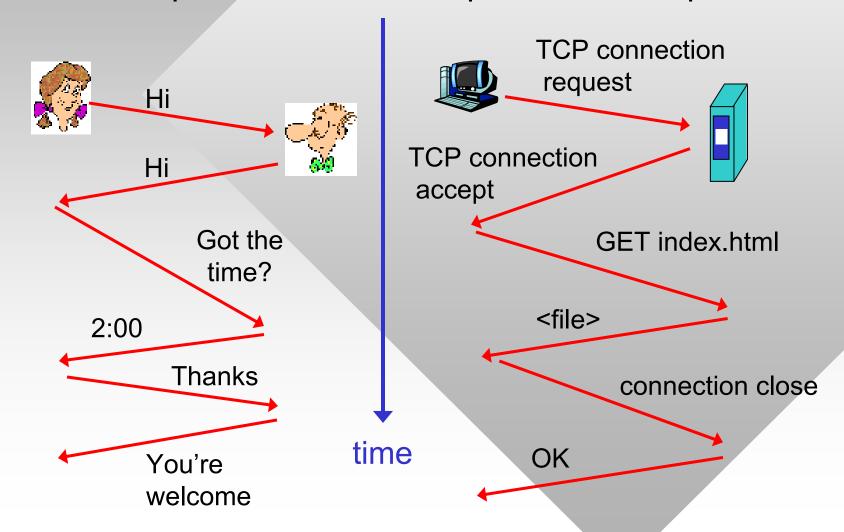
Network protocols:

- Machines rather than humans
- All communication activity in the Internet governed by protocols

Protocols define format, order of messages sent and received among network entities, and actions taken on message transmission/receipt

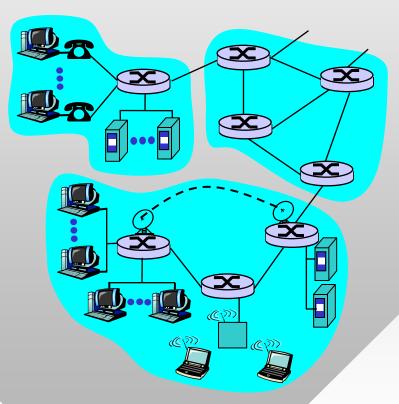
What's a Protocol?

A human protocol and a computer network protocol:



Closer Look at Network Structure

- Network edge:
 - Applications and hosts
- Network core:
 - Routers
 - Links
- How large is the edge?
 - Billions of hosts, trillions of web pages, zettabytes of information
- Large ISPs form the Internet backbone
 - Terabits per second router speed



Chapter 1: Roadmap

- 1.1 What is the Internet?
- 1.2 Network edge
- 1.3 Network core
- 1.4 Network access and physical media
- 1.5 Internet structure and ISPs
- 1.6 Delay & loss in packet-switched networks
- 1.7 Protocol layers, service models
- 1.8 History

Network Edge

The edge:

 Responsible for almost all data supply/demand

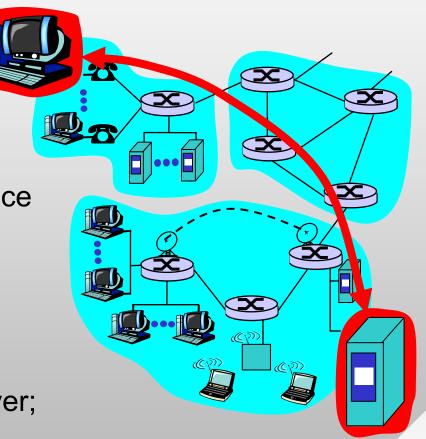
Protocols impact performance

Client/server model

- Client host requests, receives service from always-on server
- Example: web browser/server; email client/server

Peer-to-peer (P2P) model:

- Minimal use of dedicated servers; user hosts talk to each other
 - Example: BitTorrent



Network Edge: Reliable Service

- Goal: data transfer between sockets
- TCP Transmission
 Control Protocol
 - Internet's reliable service
- Connection-oriented
 - Handshaking: send connection messages (prepare) for data transfer ahead of time
 - Set up state in two communicating hosts

TCP service [RFC 793]

- Reliable, in-order bytestream data transfer
 - Packet loss handled through acknowledgements and retransmissions
- Flow control:
 - Sender won't overwhelm receiver
- Congestion control:
 - Senders reduce transmission rate when network becomes congested

Network Edge: Unreliable Service

- Goal: data transfer between sockets
 - Same as before!
- UDP User Datagram Protocol [RFC 768]:
 - Connectionless
 - Unreliable data transfer
 - No flow control
 - No congestion control
- Less overhead and delay
 - TCP connection setup & termination is 7 packets
 - TCP retransmission delay is potentially unbounded

Apps using TCP:

 HTTP (Web), FTP (file transfer), SSH (remote login), SMTP (email)

Apps using UDP:

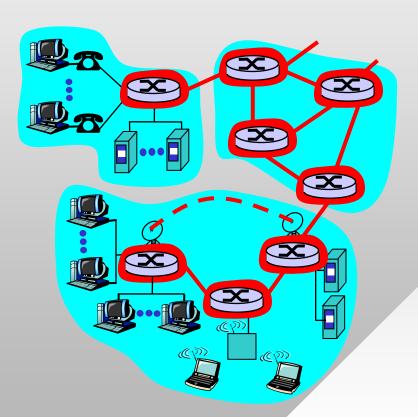
- DNS, SNMP
 - Short (single-packet) transfers
 - No need for congestion management
- Streaming media, online games, IP telephony
 - More sensitive to delay than packet loss

Chapter 1: Roadmap

- 1.1 What is the Internet?
- 1.2 Network edge
- 1.3 Network core
- 1.4 Network access and physical media
- 1.5 Internet structure and ISPs
- 1.6 Delay & loss in packet-switched networks
- 1.7 Protocol layers, service models
- 1.8 History

The Network Core

- Supports end-host communication
- Fundamental question: how is data transferred through the network?
 - Circuit switching: dedicated circuit per call (telephone network, origin 1800s)
 - Packet-switching: data sent in discrete "chunks" (1960s)
- Notation
 - Call = connection = flow



Network Core: Packet Switching

- End-to-end data stream divided into packets
 - Packets of users A and B
 share network resources
 - Each packet uses full link bandwidth

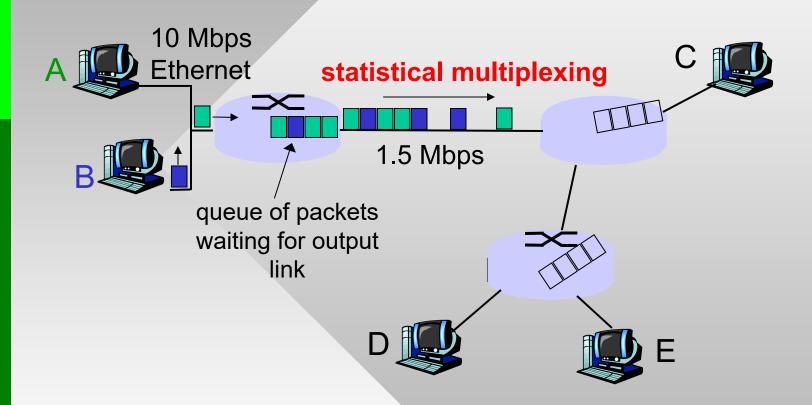
Resource contention:

- Aggregate resource demand can exceed amount available
- Congestion: packets queue, wait for link use

Store-and-forward

- Packets move one hop (router) at a time
- Node receives complete packet before forwarding

Packet Switching: Statistical Multiplexing

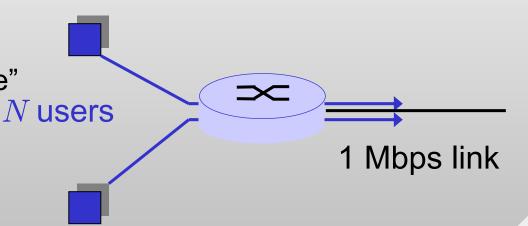


Sequence of A's and B's packets does not have a fixed pattern → statistical multiplexing

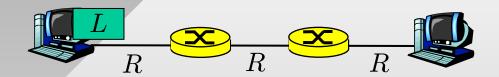
Packet Switching vs. Circuit Switching

Packet switching allows more users than circuit switching

- 1 Mbps link
- Each user:
 - 100 Kbps when "active"
 - Active 10% of time
- Circuit-switching:
 - Supports 10 users
- Packet switching:
 - With 35 users, probability that more than 10 are active is 0.0424%; with 50 users 0.94%
 - Max 100 users (if perfectly interleaved)



Packet Switching: Store-and-Forward



- Takes L/R seconds to transmit (push out) packet of L bits on to link of R bps
- Entire packet must arrive at router before it can be transmitted on next link: store and forward
- Path delay = 3L/R

Example:

- L=7.5 Mbits
- $R=1.5~\mathrm{Mbps}$
- End-to-end delay = 15 sec