<u>CSCE 463/612</u> <u>Networks and Distributed Processing</u> <u>Spring 2024</u>

Transport Layer

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February 23, 2024

Chapter 3: Transport Layer

Our goals:

- Understand principles behind transport layer services:
 - Multiplexing/demultiplexing
 - Reliable data transfer
 - Flow control
 - Congestion control
- Learn about transport layer protocols in the Internet:
 - UDP: connectionless transport
 - TCP: connection-oriented transport

Application (5) Transport (4) Network (3) Data-link (2) Physical (1)

Chapter 3: Roadmap

3.1 Transport-layer services

- 3.2 Multiplexing and demultiplexing
- 3.3 Connectionless transport: UDP
- 3.4 Principles of reliable data transfer
- 3.5 Connection-oriented transport: TCP
 - Segment structure
 - Reliable data transfer
 - Flow control
 - Connection management

3.6 Principles of congestion control3.7 TCP congestion control

Transport Services and Protocols

- Transport layer: logical communication between processes on different hosts
 - Relies on and enhances network-layer services
- Network layer: logical communication between hosts



Internet Transport-layer Protocols

- Reliable, in-order delivery: TCP
 - Congestion control
 - Flow control
 - Connection setup
- Unreliable, unordered delivery: UDP
 - No-frills extension of "besteffort" IP
- Services not available:
 - Delay or loss guarantees
 - Bandwidth guarantees



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Multiplexing/Demultiplexing



How Demultiplexing Works

- Host receives IP datagrams
 - Each datagram has source IP address and destination IP address
- Each datagram carries one transport-layer header
 - Transport header starts with source and destination port numbers
- Kernel uses port numbers to direct packets to appropriate socket or reject the message
 - Each port # is a 16-bit unsigned integer (1-65535)



Connectionless Demultiplexing

- Create a SOCK_DGRAM
 socket
- Bind the socket
 - Server: specify a well-known port (e.g., 53 for DNS)
 - Client: bind to port 0 (OS assigns next available #)
- Use sendto(), recvfrom()
- Target UDP socket is identified by a 2-tuple: (dest IP address, dest port number)

- When host receives UDP segment:
 - OS checks destination port/IP in segment
 - Directs segment to the socket with a matching combination if socket is open; rejects otherwise
- IP datagrams with different source IP addresses and/or source port numbers may be directed to the same socket!

Connectionless Demultiplexing (Cont)

SP = source port, DP = destination port



Connection-Oriented Demultiplexing

- TCP socket identified by a 4-tuple:
 - Source IP address
 - Source port number
 - Destination IP address
 - Destination port number
- Receiver host uses all four values to find appropriate socket
- <u>Clients</u>: each socket must have unique port

- <u>Servers</u>: possible to have multiple TCP sockets with same port number:
 - Each socket identified by its own 4-tuple
- Web servers have different sockets for each connecting client
 - All are on port 80
 - Non-persistent HTTP may have different socket for each request

Connection-Oriented Demultiplexing (Cont)

Web server spawns a new process per connection



SP = source port, DP = destination port; S-IP = source IP, D-IP = destination IP

Connection-Oriented Demultiplexing (Cont)

Web server spawns a new thread per connection

5775 9153 P2 **P1** P4 **P**3 5775 SP: 5775 DP: 80 S-IP: B D-IP:C SP: 9153 SP: 5775 DP: 80 DP: 80 Client client server S-IP: A S-IP: B IP:B IP: A IP: C D-IP:C D-IP:C

port 80

SP = source port, DP = destination port; S-IP = source IP, D-IP = destination IP