# JetMax: Scalable Max-Min Congestion Control for High-Speed Heterogeneous Networks

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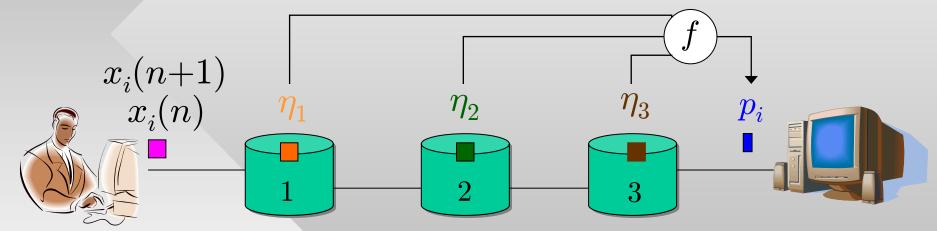
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#### <u>Agenda</u>

- Introduction
  - Explicit congestion control and its designed properties
- Analysis of existing max-min methods
  - XCP, MKC, and MKC-AVQ
  - JetMax
    - Stability and fairness
    - Performance and simulations
  - Wrap-up

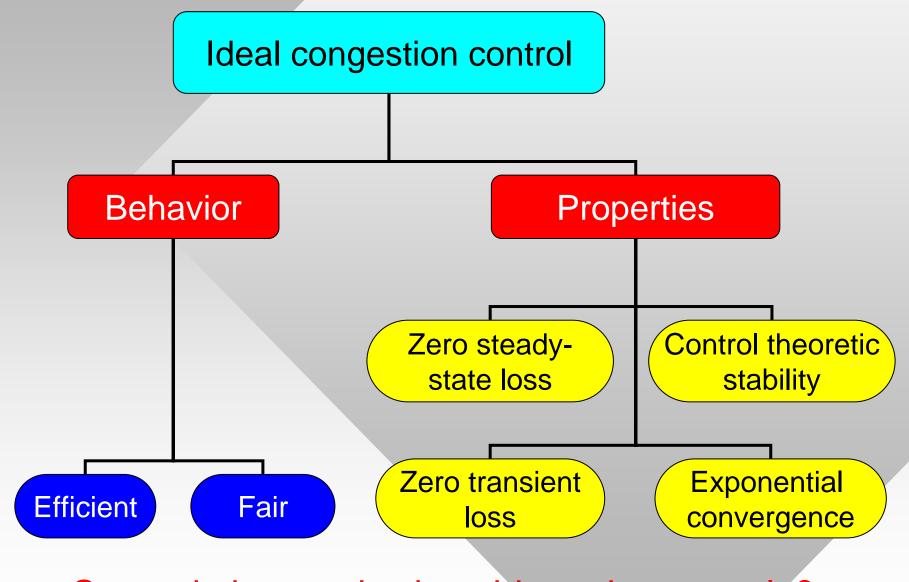
#### **Explicit Congestion Control**

- In explicit congestion control, each user i forms a control loop around its sending rate  $x_i(n)$  and feedback  $p_i(n)$ 



- Two directions emerged
  - $p_i = \Sigma \eta_i$ : additive feedback (e.g., proportional fairness)
  - $p_i = \max(\eta_l)$ : max-min feedback (e.g., max-min fairness)
- What properties should a congestion control protocol ideally possess?

#### **Ideal Congestion Control**



Can existing methods achieve these goals?

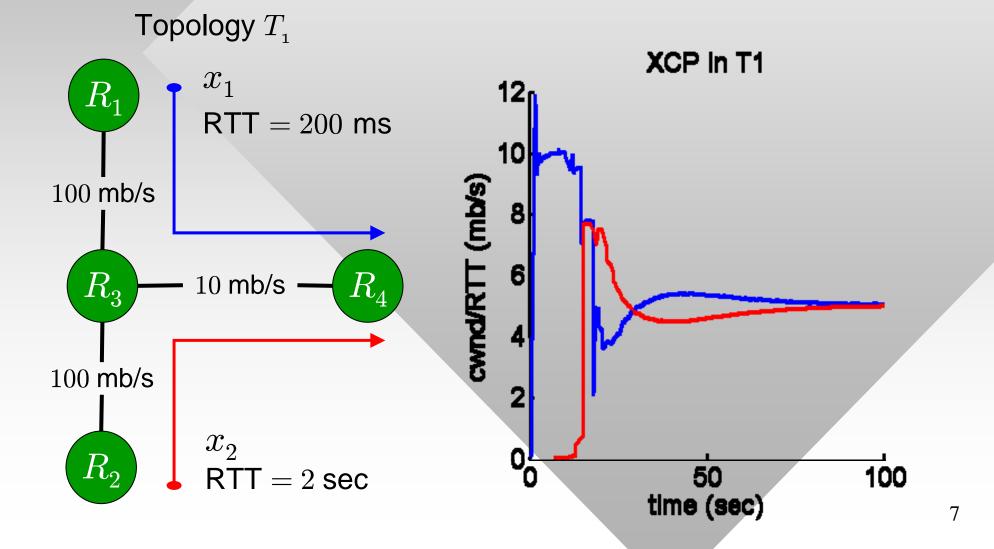
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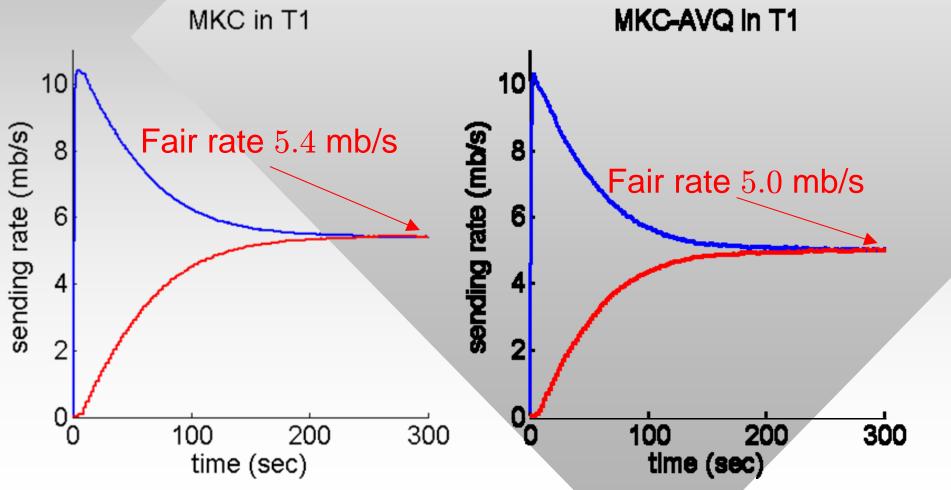
## **Existing Max-min Congestion Controls**

- We study three explicit congestion control methods
- XCP (eXplicit Control Protocol)
  - Each router implements a fairness and efficiency controller
- MKC (Max-min Kelly Control)
  - Modification of Kelly's equations for max-min feedback
- MKC-AVQ
  - Combination of MKC and Adaptive Virtual Queue (AVQ)
  - Eliminates steady-state loss of MKC

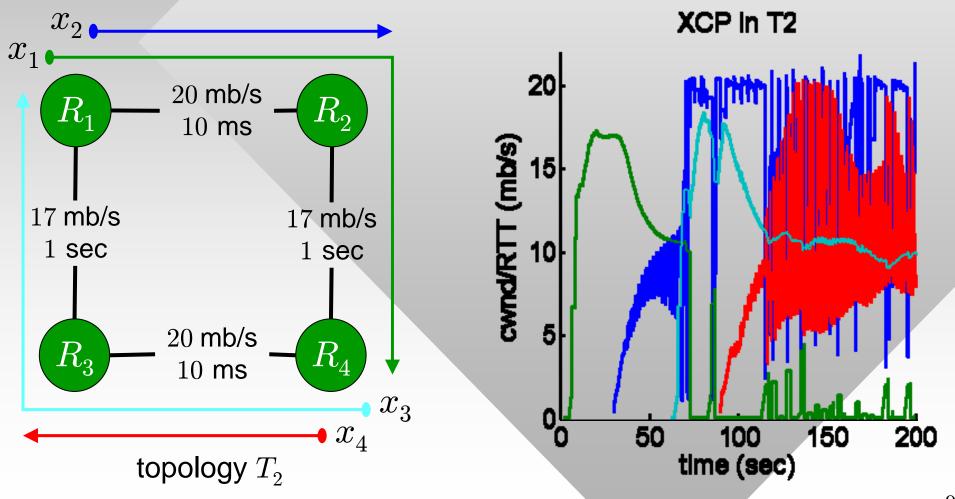
- All three should be stable in single-link networks
  - We next compare their behavior under heterogeneous delay



 MKC-based methods have no stability problems, but are very slow

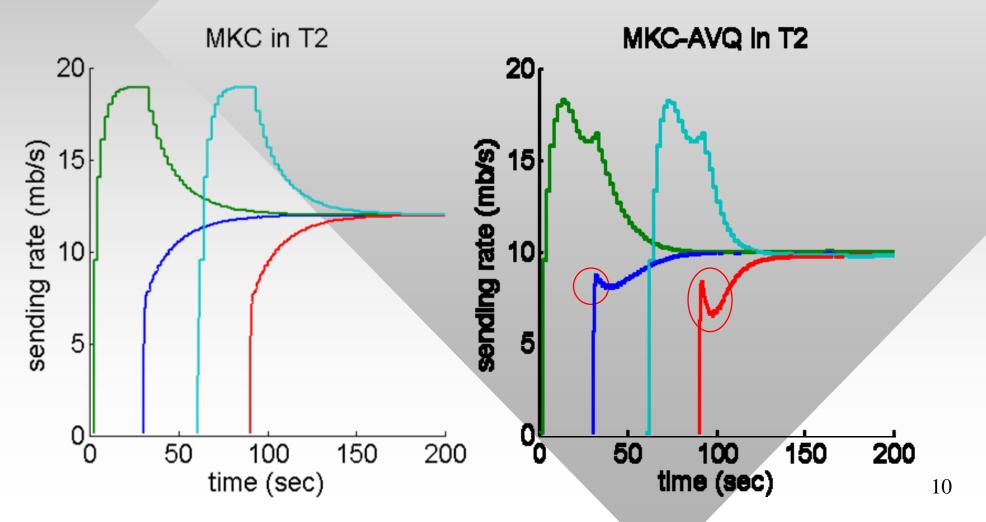


- Multi-link behavior
  - XCP example shows that bottleneck oscillations are possible



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- MKC and MKC-AVQ remain stable
  - AVQ overshoots available bandwidth in transient states, but is stable otherwise



## **Properties of Existing Max-min Methods**

- XCP
  - Unstable in certain multi-link topologies
  - Convergence rate to fairness is unknown, but can be as large as 100s of RTTs if feedback delay is heterogeneous
- MKC
  - Steady-state packet loss proportional to N
  - Fairness is reached in  $\Theta(C)$  steps
- MKC-AVQ
  - Transient overshoot proportional to N
  - Convergence rate the same as in MKC
- The bottom line current explicit-feedback methods do not satisfy our design criteria and offer little benefit over TCP

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## JetMax Design

- Assume that N<sub>l</sub>(n) is the number of users congested by router l at time n (we call these flows responsive) and w<sub>l</sub>(n) is their combined rate
  - Label the remaining flows passing through l unresponsive and assume  $u_l(n)$  is their total rate
- The main idea of JetMax is very simple
  - Divide the available bandwidth of l equally among  $N_l$  flows

 $u_l(n$ 

capacity of link l

current fair share at router *l* 

 $g_l(n) = \gamma_l$ 

combined rate of unresponsive flows

desired link utilization

number of responsive flows

## JetMax Design 2

Users receive feedback  $g_l(n)$  and utilize it in their controller RTT of user *i* 

constant

$$x_{i}(n) = x_{i}(n - D_{i}) - \tau \left( x_{i}(n - D_{i}) - g_{l}(n - D_{i}) \right)$$
  
constant delay from the bottleneck to user *i*

- The final issue is bottleneck switching
  - Routers decide to make a switch based on the virtual packet loss contained in the header:

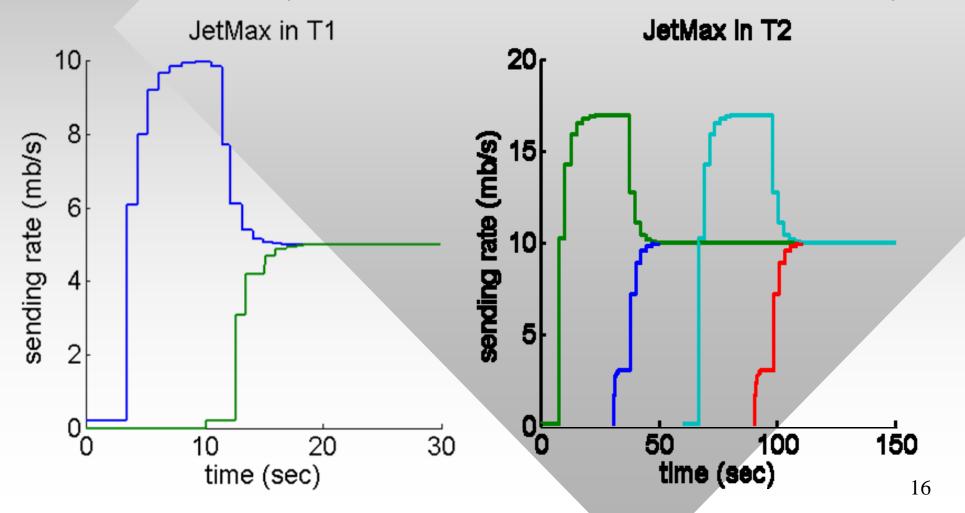
$$p_{l}(n) = \frac{w_{l}(n) + u_{l}(n) - \gamma_{l}C_{l}}{w_{l}(n) + u_{l}(n)}$$

## **JetMax Properties**

- <u>Theorem 1:</u> Under any fixed bottleneck assignment under max-min feedback, JetMax is globally asymptotically stable regardless of delay if and only if  $0 < \tau < 2$ 
  - For values of  $\tau \leq 1$ , the controller is also monotonic
- <u>Theorem 2</u>: Stationary resource allocation of JetMax is max-min fair
- <u>Theorem 3</u>: On a single link of any capacity, JetMax converges to within ε-percent of efficiency and fairness in a fixed number of RTT steps
  - For  $\tau = 0.5$  and  $\varepsilon = 1\%$ , this is 6 steps

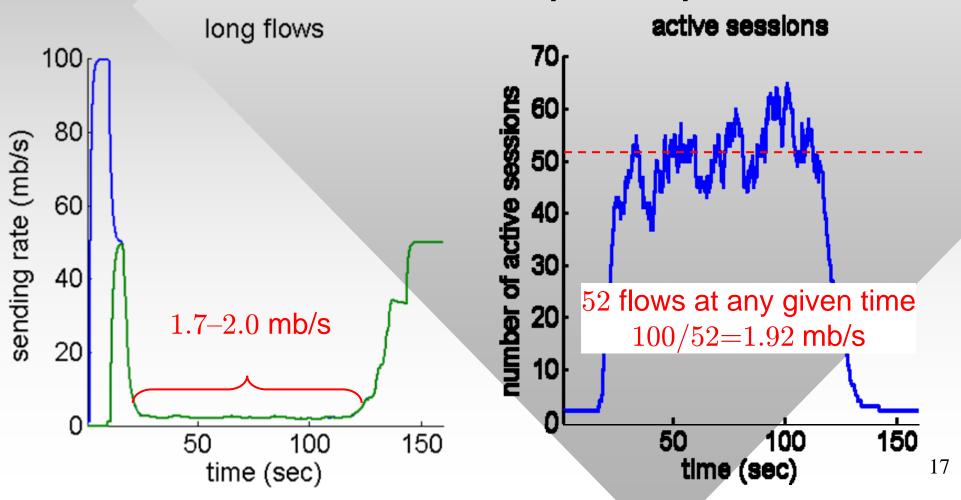
## **JetMax Simulations**

- We start with ns2 simulations in  $T_1$  and  $T_2$ 
  - Topology  $T_1$  is changed to introduce random time-varying feedback delay in the control loop no effect on stability



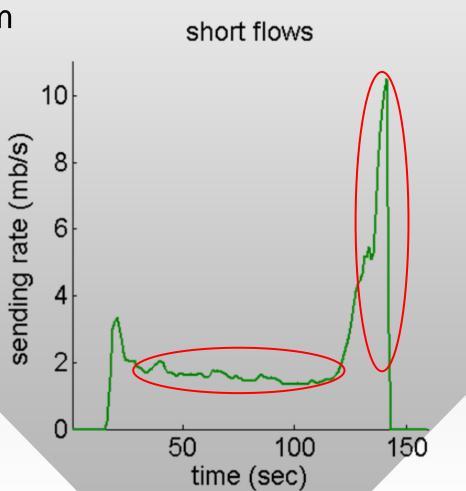
## **JetMax Simulations 2**

- Performance under mice cross-traffic
  - A single link of capacity 100 mb/s shared by two long flows and a total of 500 short flows with random RTT in [40,1040] ms and uniform packet size in [800,1300] bytes



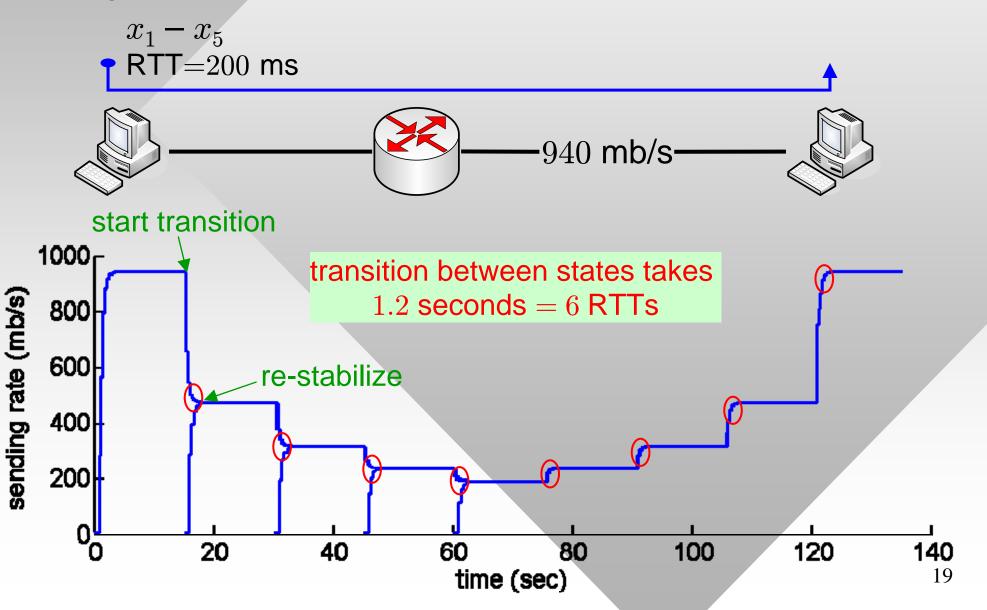
## **JetMax Simulations 3**

- Short flows
  - Each lasts for a random duration with mean 10 seconds
  - Achieve rates very close to fair-share 1.92 mb/s
- As flows depart after 120 seconds, their bandwidth is consumed by remaining flows



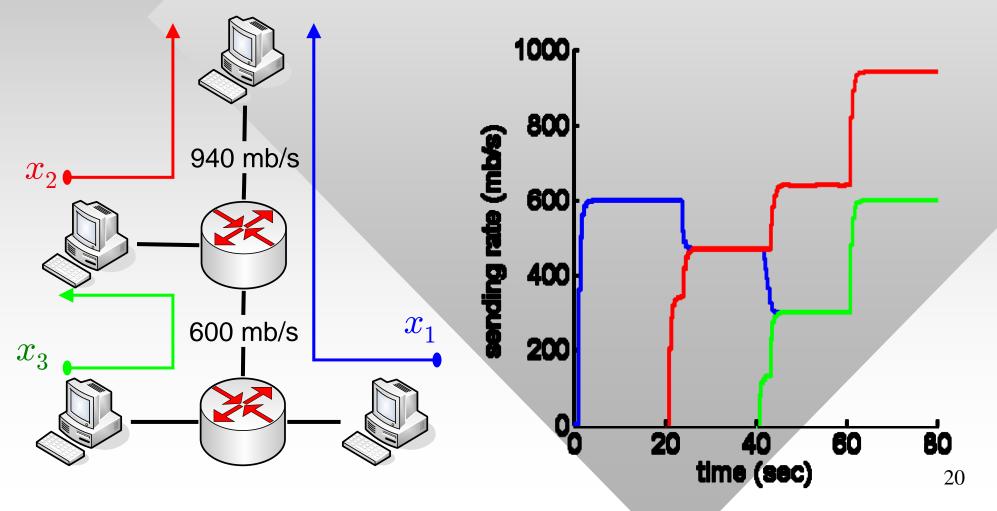
#### JetMax Linux Experiments 1

Single-link scenario



#### JetMax Linux Experiments 2

- Multi-link scenario
  - Flow  $x_1$  starts first,  $x_2$  arrives at time 20 seconds,  $x_3$  arrives at 40 seconds, and  $x_1$  departs at 60 seconds



#### Wrap-up

- JetMax summary
  - Converges to the stationary state in the same number of RTTs regardless of link capacity and the number of flows
  - Achieves tunable utilization and zero loss
  - Is monotonic and stable under any time-varying delay
- More in the paper
  - Estimation of responsive/unresponsive rates
  - Bottleneck membership maintenance
  - Avoidance of transient overshoot
  - Additional simulations and experiments
- Linux and ns2 code
  - http://irl.cs.tamu.edu/mkc